

Answers to the Exercises -- Chapter 2

SECTION 1

Exercises 1 and 2 answered together:

- a. Not a sentence
- b. Informal notation

$$\begin{array}{c} \sim Q \leftrightarrow \sim R \\ \text{F} \\ \wedge \\ \sim Q \quad \sim R \\ \text{F} \quad \text{T} \\ | \quad | \\ Q \quad R \\ \text{T} \quad \text{F} \end{array}$$

- c. Official notation

$$\begin{array}{c} \sim(Q \leftrightarrow R) \\ \text{T} \\ | \\ Q \leftrightarrow R \\ \text{F} \\ \wedge \\ Q \quad R \\ \text{T} \quad \text{F} \end{array}$$

- d. Not a sentence
- e. Informal notation

$$\begin{array}{c} (P \rightarrow Q) \vee (R \rightarrow \sim Q) \\ \text{T} \\ \wedge \\ P \rightarrow Q \quad R \rightarrow \sim Q \\ \wedge \quad \wedge \\ P \quad Q \quad R \quad \sim Q \\ \text{F} \quad | \\ Q \end{array}$$

- f. Not a sentence
- g. Informal notation

$$\begin{array}{c} P \wedge Q \rightarrow (Q \rightarrow R \vee Q) \\ \text{T} \\ \wedge \\ P \wedge Q \quad Q \rightarrow R \vee Q \\ \wedge \quad \wedge \\ P \quad Q \quad Q \quad R \vee Q \\ \text{F} \quad | \\ R \quad Q \\ \text{T} \end{array}$$

h. Informal notation

$$\begin{array}{c}
 P \leftrightarrow (P \leftrightarrow Q \wedge R) \\
 | \\
 F \\
 \wedge \\
 | \\
 P \quad P \leftrightarrow Q \wedge R \\
 | \quad | \\
 T \quad F \\
 \wedge \\
 | \\
 P \quad Q \wedge R \\
 | \quad | \\
 T \quad F \\
 \wedge \\
 | \\
 Q \quad R \\
 | \\
 F
 \end{array}$$

i. Informal notation

$$\begin{array}{c}
 P \vee (Q \rightarrow P) \\
 | \\
 T \\
 \wedge \\
 | \\
 P \quad Q \rightarrow P \\
 | \quad | \\
 T \quad \wedge \\
 | \quad Q \quad P
 \end{array}$$

SECTION 2

- 1. a. $R \wedge P$
- b. $W \vee R$
- c. $\sim R \wedge T$
- d. $R \wedge S$
- e. $Q \leftrightarrow R$

2 and 3 answered together

- a. $S \vee V$ false
- b. $R \leftrightarrow S$ true
- c. $R \wedge S$ false
- d. $Q \vee T$ false
- e. $Q \wedge S$ false

SECTION 3

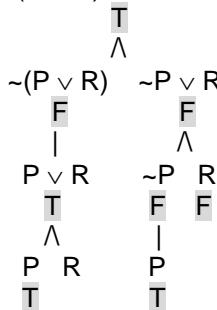
1. a. $\sim(P \vee (Q \wedge R))$

$$\begin{array}{c}
 | \\
 F \\
 | \\
 P \vee (Q \wedge R) \\
 | \\
 T \\
 \wedge \\
 | \\
 P \quad Q \wedge R \\
 | \quad | \\
 T \quad \wedge \\
 | \quad Q \quad R
 \end{array}$$

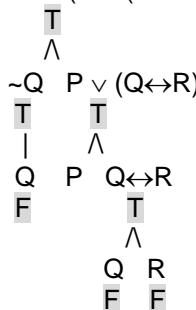
b. $\sim P \vee (Q \wedge R)$

$$\begin{array}{c}
 | \\
 F \\
 \wedge \\
 | \\
 \sim P \quad Q \wedge R \\
 | \quad | \\
 F \quad F \\
 | \quad \wedge \\
 P \quad Q \quad R \\
 | \quad | \quad | \\
 T \quad F \quad F
 \end{array}$$

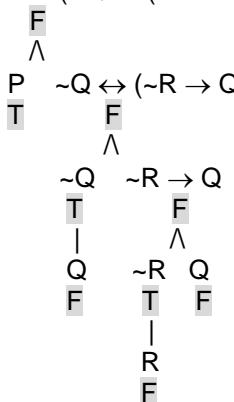
c. $\neg(P \vee R) \leftrightarrow \neg P \vee \neg R$



d. $\neg Q \wedge (P \vee (Q \leftrightarrow R))$



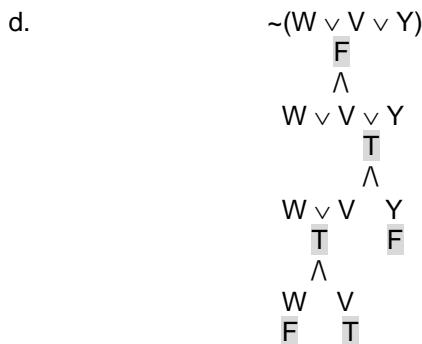
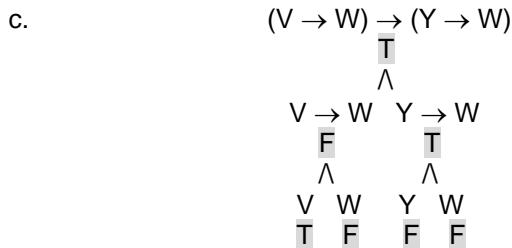
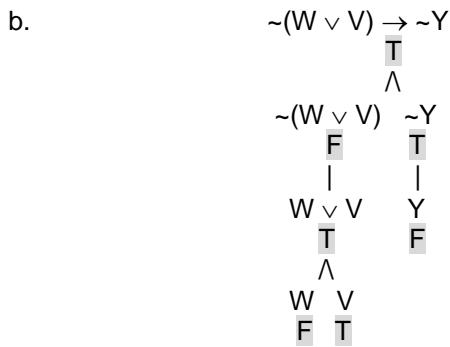
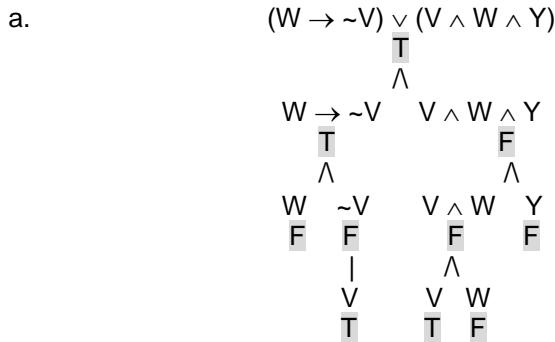
e. $P \rightarrow (\neg Q \leftrightarrow (\neg R \rightarrow Q))$



2. a. $\neg V \leftrightarrow \neg W$; "won't" is a negation with narrow scope.
b. $\neg V \rightarrow (Y \rightarrow W \wedge V)$; "both" gives rise to a conjunction with narrow scope since it splits the names from the predicate. The comma prevents Y and V from occurring together.
c. $Y \vee (W \vee V)$; "unless" is a disjunction sign.
d. $(Y \wedge \neg V) \vee (V \wedge \neg W)$
3. a. Only if Veronica doesn't leave will William leave, or Veronica and William and Yolanda will all leave.
 $(\text{Only if Veronica doesn't leave will William leave}) \vee (\text{Veronica and William and Yolanda will leave})$
 $(\text{William will leave} \rightarrow \text{Veronica doesn't leave}) \vee (V \wedge W \wedge Y)$
 $(W \rightarrow \neg V) \vee (V \wedge W \wedge Y)$
- b. If neither William nor Veronica leaves, Yolanda won't either
If neither William [leaves] nor Veronica leaves, [then] Yolanda won't [leave]
 $\neg(W \vee V) \rightarrow \neg Y$
- c. If William will leave if Veronica leaves, then he will surely leave if Yolanda leaves
If (William will leave if Veronica leaves) then ([William] will leave if Yolanda leaves)
 $(V \rightarrow W) \rightarrow (Y \rightarrow W)$

- d. Neither William nor Veronica nor Yolanda will leave
 $\sim(W \vee V \vee Y)$

4. "Veronica leaves but neither William nor Yolanda leaves" corresponds to the truth-value assignment: V --- true; W --- false; Y --- false. We use parse trees to compute the truth values of the complex sentences.



5. a. Sally will run and win unless she quits
 $(Sally \text{ will run and } [Sally \text{ will}] \text{ win}) \vee ([Sally] \text{ quits})$
 $(R \wedge W) \vee Q$
- b. Sally will win exactly in case she runs without quitting
Sally will win exactly in case (she runs [and doesn't] quit)
 $W \leftrightarrow (R \wedge \neg Q)$
- c. Sally, who will run, will win if she doesn't quit
Sally will run, and Sally will win if she doesn't quit
 $R \wedge (\neg Q \rightarrow W)$
- d. Sally will run and quit, but she will win anyway
Sally will run and quit, and she will win
 $(R \wedge Q) \wedge W$

SECTION 4

1. a. None; if we had $\neg\neg Q$ instead of Q it would be an instance of MTP.
- b. Simplification
- c. Double Negation
- d. MTP
- e. CB
- f. None.
- g. BC
- h. Adjunction
- i. None
2. a. $\neg W \leftrightarrow \neg X$ by CB; also $\neg X \leftrightarrow \neg W$ by CB
- b. $\neg\neg W$ by MTP
- c. Nothing
- d. $\neg W$ by S; also $\neg X$ by S
- e. $W \rightarrow \neg X$ by BC; also $\neg X \rightarrow W$ by BC
- f. Nothing

SECTION 5 Derivations of numbered theorems not given

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SECTION 7

1. a. All fine
- b. In line 8, the sentence that can be inferred from 7 by RT39 is $W \rightarrow \neg S$.
- 2, 3, 4, 5: Derivations of numbered theorems not given

SECTION 8

1. a. All fine
- b. Line 4: MTP does not apply;
Line 8: BC (biconditional to conditional) does not apply; we could use CB;
Line 11: MP does not apply to biconditionals; you have to split the biconditional into conditionals first using BC.
- c. Line 2: the result of applying DM to pr2 is $\neg Y \wedge \neg\neg Z$ rather than $\neg Y \wedge Z$.
Line 3: NC doesn't apply; the NC would generate line 3 if line 2 were $Y \wedge \neg Z$.
Line 4: Line 4 is not available at line 4; it may not be cited to justify itself. The sentence could be generated by applying MT to line 3 and pr1.

2. a. $U \wedge V \rightarrow X$ <use dm>

$\sim V \rightarrow Y$

$X \vee Y \rightarrow Z$

$\therefore \sim Z \rightarrow \sim U$

1. Show $\sim Z \rightarrow \sim U$

2.	$\sim Z$	ass cd
3.	$\sim(X \vee Y)$	pr3 2 mt
4.	$\sim X \wedge \sim Y$	3 dm
5.	$\sim X$	4 s
6.	$\sim Y$	4 s
7.	$\sim(U \wedge V)$	pr1 5 mt
8.	$\sim U \vee \sim V$	7 dm
9.	$\sim\sim V$	6 pr2 mt
10.	$\sim U$	8 9 mtp cd

b. $(X \rightarrow Y) \rightarrow Z$ <use nc>

$\sim Z$

$V \rightarrow Y$

$\therefore \sim V$

1. Show $\sim V$

2.	$\sim(X \rightarrow Y)$	pr1 pr2 mt
3.	$X \wedge \sim Y$	2 nc
4.	$\sim Y$	3 s
5.	$\sim V$	4 pr3 mt dd

c. $P \vee Q$

$Q \rightarrow S$

$U \vee \sim S$

$P \vee S \rightarrow R$

$R \rightarrow U$

$\therefore U$

1. Show U

2.	Show $P \rightarrow U$	
3.	P	ass cd
4.	$P \vee S$	3 add
5.	R	4 pr4 mp
6.	U	5 pr5 mp cd
7.	Show $Q \rightarrow U$	
8.	Q	ass cd
9.	S	8 pr2 mp
10.	$\sim\sim S$	9 dn
11.	U	10 pr3 mtp cd
12.	U	pr1 2 7 sc
13.		12 dd

SECTION 9

1. a. $\neg(P \leftrightarrow Q)$ <use nb>
 $R \vee P$
 $\neg Q \rightarrow R$
 $\therefore R$

1. Show R

2.	$\neg R$	ass id
3.	P	2 pr2 mtp
4.	$P \leftrightarrow \neg Q$	pr1 nb
5.	$P \rightarrow \neg Q$	4 bc
6.	$\neg Q$	3 5 mp
7.	R	6 pr3 mp
8.		2 7 id

- b. $W \rightarrow U$ <use cdj>
 $\neg W \rightarrow V$
 $\therefore U \vee V$

1. Show $U \vee V$

Show $\neg U \rightarrow V$		
3.	$\neg U$	ass cd
4.	$\neg W$	3 pr1 mt
5.	V	4 pr2 mp cd
6.	$U \vee V$	2 cdj dd

- c. $P \vee (Q \wedge S)$
 $R \vee Q$
 $S \vee \neg P$
 $Q \rightarrow \neg S$
 $\therefore R$

1. Show R

2.	$\neg R$	ass id
3.	Q	2 pr2 mtp
4.	$\neg S$	3 pr4 mp
5.	$\neg Q \vee \neg S$	4 add
6.	$\neg(Q \wedge S)$	5 dm
7.	P	6 pr1 mtp
8.	$\neg\neg P$	7 dn
9.	S	8 pr3 mtp
10.		4 9 id

SECTION 10

1. a. $(R \leftrightarrow S) \vee (R \leftrightarrow \neg S)$; tautology

R	S	$(R \leftrightarrow S) \vee (R \leftrightarrow \neg S)$
T	T	T
T	F	T
F	T	T
F	F	T

b. $R \leftrightarrow (S \leftrightarrow R)$; not a tautology

R	S	$R \leftrightarrow (S \leftrightarrow R)$
T	F	F

c. $R \vee (S \wedge T) \rightarrow R \wedge (S \vee T)$; not a tautology

R	S	T	$R \vee (S \wedge T) \rightarrow R \wedge (S \vee T)$
F	T	T	F

d. $\neg U \rightarrow (U \rightarrow \neg V)$; tautology

U	V	$\neg U \rightarrow (U \rightarrow \neg V)$
T	T	T
T	F	T
F	T	T
F	F	T

e. $(\neg R \leftrightarrow R) \rightarrow S$; tautology

R	S	$(\neg R \leftrightarrow R) \rightarrow S$
T	T	T
T	F	T
F	T	T
F	F	T

f. $(S \wedge T) \vee (S \wedge \neg T) \vee \neg S$; tautology

T	S	$(S \wedge T) \vee (S \wedge \neg T) \vee \neg S$
T	T	T
T	F	T
F	T	T
F	F	T

SECTION 11

a. $\begin{aligned} U \wedge V &\rightarrow X && \text{NO} \\ \neg V &\rightarrow U \\ X \vee V &\rightarrow U \\ \therefore V &\rightarrow \neg U \end{aligned}$

U	V	X	$U \wedge V \rightarrow X$	$\neg V \rightarrow U$	$X \vee V \rightarrow U$	$V \rightarrow \neg U$
T	T	T	T	T	T	F

b. $(X \rightarrow Y) \rightarrow Z$ YES

$$\begin{array}{l} \sim Z \\ \therefore \sim Y \end{array}$$

X	Y	Z	$(X \rightarrow Y) \rightarrow Z$	$\sim Z$	$\sim Y$
T	T	T	T	F	F
T	T	F	F	T	F
T	F	T	T	F	T
T	F	F	T	T	T
F	T	T	T	F	F
F	T	F	F	T	F
F	F	T	T	F	T
F	F	F	F	T	T

c. $\sim(P \leftrightarrow Q)$ YES

$$\begin{array}{l} R \vee P \\ \sim Q \rightarrow R \\ \therefore R \end{array}$$

P	Q	R	$\sim(P \leftrightarrow Q)$	$R \vee P$	$\sim Q \rightarrow R$	R
T	T	T	F	T	T	T
T	T	F	F	T	T	F
T	F	T	T	T	T	T
T	F	F	T	T	F	F
F	T	T	T	T	T	T
F	T	F	T	F	T	F
F	F	T	F	T	T	T
F	F	F	F	F	F	F

d. $S \vee T$ NO

$$\begin{array}{l} W \vee S \\ \sim T \vee \sim S \\ \therefore \sim S \end{array}$$

S	T	W	$S \vee T$	$W \vee S$	$\sim T \vee \sim S$	$\sim S$
T	F	T	T	T	T	F

e. $W \rightarrow U$ YES

$$\begin{array}{l} \sim W \rightarrow V \\ \therefore U \vee V \end{array}$$

U	V	W	$W \rightarrow U$	$\sim W \rightarrow V$	$U \vee V$
T	T	T	T	T	T
T	T	F	T	T	T
T	F	T	T	T	T
T	F	F	T	F	T
F	T	T	F	T	T
F	T	F	T	T	T
F	F	T	F	T	F
F	F	F	T	F	F

f. $P \leftrightarrow \sim Q$ NO

$$\begin{array}{l} Q \rightarrow R \vee P \\ R \rightarrow \sim Q \vee \sim P \\ \therefore Q \vee R \end{array}$$

P	Q	R	$P \leftrightarrow \sim Q$	$Q \rightarrow R \vee P$	$R \rightarrow \sim Q \vee \sim P$	$Q \vee R$
T	F	F	T	T	T	F

g. $P \vee (Q \wedge S)$ YES
 $S \vee Q$
 $S \vee \neg P$
 $\therefore S$

P	Q	S	$P \vee (Q \wedge S)$	$S \vee Q$	$S \vee \neg P$	S
T	T	T	T	T	T	T
T	T	F	T	T	F	F
T	F	T	T	T	T	T
T	F	F	T	F	F	F
F	T	T	T	T	T	T
F	T	F	F	T	T	F
F	F	T	F	T	T	T
F	F	F	F	F	T	F

h. $P \wedge (Q \vee S)$ NO
 $S \vee Q$
 $S \vee P$
 $\therefore S$

P	Q	S	$P \wedge (Q \vee S)$	$S \vee Q$	$S \vee P$	S
T	T	F	T	T	T	F